

BBRP22 based security technique for data security

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Abstract. Nowadays data plays major role in the current internet world. People generate more data through online media. The generated data are not secured properly and can easily get hacked by the hackers. To overcome from this data security issue, we have planned to apply BBRP25 technique to secure the data. BBRP22 technique has 6 phases. To apply this technique, we must follow the given below steps: 1. To track down the great key S; 2. To get the Z1 & Z2 values from indivisible numbers; 3. To view as the an and b values in lattice; to find the values with the assistance k; and trade x and y values. 4. To get the T-test values and match those numbers. 5. To find the n esteem by utilizing k. 6. Apply the n values in legitimate condition; and trade x and y values. The data will be fully secured by applying BBRP22 strategy while compared to ChaCha.

Keywords: BBRP22, encryption, ChaCha, security.

1. Introduction

Nowadays data plays major role in the current internet world. People generate more data through online media. The generated data are not secured properly and can easily get hacked by the hackers. To overcome from this data security issue, we have planned to apply BBRP22 technique to secure the data. ChaCha family has introduced the first ChaCha20/20. ChaCha20/4 mainly used for the reduced the encryption time has comparing with ChaCha20. ChaCha20 technique has mainly using for speed of the encryption but “not increasing the security of current data”. To “beat this disadvantage of ChaCha20/4”; presented the novel strategy BBRP (Bagath Basha and RajaPrakash) 22. The time complexity has improved by using ChaCha [1]. They studied about the 3Vs [2] ChaCha technique has proposed by “Bernstein D. J.” [15]. The attacked the ChaCha using XOR method [3]. To create new idea using hash method [4]. To produced the novel attack for ChaCha [5] and Double A [6]. To created new calculation method for security [7]. SRB18 system is utilized to give security to data [8]. SRB21-1 and SRB21-2 technique are utilized to provide the data security [9],[10]. CBB21 [11], CBB22 [12], CBB20 [13], and RJB25 methods are used to provide security of data [14-17]. They proposed 7 phases for giving the security of the information [18].

2. Methodology

Nowadays data plays major role in the current internet world. People generate more data through online media. The generated data are not secured properly and can easily get hacked by the hackers. To overcome from this data security issue, we have planned to apply BBRP25 technique to secure the data. BBRP22 technique has 6 phases. To apply this technique, we must follow the given below steps: 1. To track down the great key S; 2. To get the Z1 & Z2 values from indivisible numbers; 3. To view as the an and b values in lattice; to find the values with the assistance k; and trade x and y values. 4. To get the T-test values and match those numbers. 5. To find the n esteem by utilizing k. 6. Apply the n values in legitimate condition; and trade x and y values.

3. BBRP22 encryption technique

1=> Fetch the input matrix (IP) form social media.

2=> To get the key from IP.

3=> Make it two parts using equation (1), (2), and (3).

4=> T-Test Formula =

$$\frac{(\bar{Z}_1 - \bar{Z}_2)}{\sqrt{\left(\frac{A_1^2}{N_1}\right) + \left(\frac{A_2^2}{N_2}\right)}} \quad (1)$$

5=> Get the values from using equation (1)

$$6=> \quad \bar{Z}_1 = \sum \frac{Z_1}{N_1} \quad \bar{Z}_2 = \sum \frac{Z_2}{N_2} \quad (2)$$

$$7=> A_1 = \sqrt{\frac{\sum (Z_1 - \bar{Z}_1)^2}{(N_1 - 1)}} \quad A_2 = \sqrt{\frac{\sum (Z_2 - \bar{Z}_2)^2}{(N_2 - 1)}} \quad (3)$$

8=> Get pair values from left to right.

9=> Get n and k value.

$$10=> x_n + y_n = (x + y) (x_{n-1} - x_{n-2}y + x_{n-3}y^2 + \dots + x_{n-2x} - y_{n-1}) \quad (4)$$

4. Result and discussion

- The BBRP22 technique is developed from RJB25 technique [20].

$$IP = \begin{bmatrix} 301/3 & 302/3 & 303/3 & 304/3 & 305/3 \\ 306/3 & 307/3 & 308/3 & 309/3 & 310/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Where IP (Input of the matrix).

- Prime Numbers – 1, 3, 5, 7, 11, 13, 17, 19, 23
 - X₁ = 3, 5, 7, 11
 - X₂ = 13, 17, 19, 23

Table 1. Z1 and Z2 values.

Z ₁	(Z ₁ – \bar{Z}_1)	(Z ₁ – \bar{Z}_1) ²	Z ₂	(Z ₂ – \bar{Z}_2)	(Z ₂ – \bar{Z}_2) ²
3	-3.5	12.25	13	-5	25
5	-1.5	2.25	17	-1	1

Table 1. (continued).

7	0.5	0.25	19	1	1
11	4.5	20.25	23	5	25
$\sum (X_1 - \bar{X}_1)^2$		35	$\sum (X_2 - \bar{X}_2)^2$		52

- **Using Equation (2) and (7) and Table 1**
 - $\bar{Z}_1=6.5$
 - $\bar{Z}_2=18$
- **Using Equation (3) and (6)**
- $Z_1 = \sqrt{(35/(4-1))}$ $Z_2 = \sqrt{(52/(4-1))}$
- $Z_1 = \sqrt{(35/3)}$ $Z_2 = \sqrt{(52/3)}$
- $S_1 = 11.66$ $S_2 = 17.33$
- **Using Equation (1) and (8)**
- $TTF = (6.5 - 18)/\sqrt{((11.66^2/4) + (17.33^2/4))}$
- $TTF = 11.5/\sqrt{109.06}$

Step 1: (1,1)

$$TTE = \begin{bmatrix} 301/3 & 302/3 & 303/3 & 304/3 & 305/3 \\ 306/3 & 307/3 & 308/3 & 309/3 & 310/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Where TTE is T-Test Encryption

Step 2: (5,1)

$$TTE = \begin{bmatrix} 301/3 & 306/3 & 303/3 & 304/3 & 305/3 \\ 302/3 & 307/3 & 308/3 & 309/3 & 310/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Step 3: (0,9)

$$TTE = \begin{bmatrix} 310/3 & 306/3 & 303/3 & 304/3 & 305/3 \\ 302/3 & 307/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Step 4: (0,6)

$$TTE = \begin{bmatrix} 307/3 & 306/3 & 303/3 & 304/3 & 305/3 \\ 302/3 & 310/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

- To find the n value; a=2,b=3, k=2,n=5, **n is odd number**
- **Using Equation (4) and (5)**

- $x^2+y^3 = (2+3)(2^{5-1}-2^{5-2}3+2^{5-3}3^2-2^{5-4}3^3-2^{5-5}3^4+3^{5-4}2-3^{5-3}2+3^{5-2}2-3^{5-1})$
- $x^2+y^3 = (5)(2^4-2^33+2^29-2^127-2^081+3^12-3^22+3^32-3^4)$
- $x^2+y^3 = (5)(16-24+36-54-81+6-18+54-81)$
- $x^2+y^3 = (5)(-146)$
- $x^2+y^3 = (5,1), (4,6)$

Step 5: (5,1)

$$TTOE = \begin{bmatrix} 307/3 & 302/3 & 303/3 & 304/3 & 305/3 \\ 306/3 & 310/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Where TTOE is T-Test Odd Encryption

Step 6: (4, 6)

$$TTOE = \begin{bmatrix} 307/3 & 302/3 & 303/3 & 304/3 & 310/3 \\ 306/3 & 305/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Decryption

- Pair the odd value from right left (6, 4) and (1, 5) and swap it those numbers.

Step 1: (6, 4)

$$TTOD = \begin{bmatrix} 307/3 & 302/3 & 303/3 & 304/3 & 305/3 \\ 306/3 & 310/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Where TTOD is T-Test Odd Decryption

Step 2: (1, 5)

$$TTE = \begin{bmatrix} 307/3 & 306/3 & 303/3 & 304/3 & 305/3 \\ 302/3 & 310/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

- Pair “the T-test value from right to left (6,0), (9,0), (1,5), and (1,1) and swap it those numbers”.

Step 3: (6,0)

$$TTD = \begin{bmatrix} 310/3 & 306/3 & 303/3 & 304/3 & 305/3 \\ 302/3 & 307/3 & 308/3 & 309/3 & 301/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Where TTE is T-Test Decryption

Step 4: (9,0)

$$TTD = \begin{bmatrix} 301/3 & 306/3 & 303/3 & 304/3 & 305/3 \\ 302/3 & 307/3 & 308/3 & 309/3 & 310/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Step 5: (1,5)

$$TTD = \begin{bmatrix} 301/3 & 302/3 & 303/3 & 304/3 & 305/3 \\ 306/3 & 307/3 & 308/3 & 309/3 & 310/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

Step 6: (1,1)

$$TTD = \begin{bmatrix} 301/3 & 302/3 & 303/3 & 304/3 & 305/3 \\ 306/3 & 307/3 & 308/3 & 309/3 & 310/3 \\ 311/3 & 312/3 & 313/3 & 314/3 & 315/3 \\ 316/3 & 317/3 & 318/3 & 319/3 & 320/3 \\ 321/3 & 322/3 & 323/3 & 324/3 & 325/3 \end{bmatrix}$$

The proposed algorithm BBRP22 compare the performance with existing method “ChaCha”. The existing method is to do the process for move all diagonal values into the first column. The three by three matrix has (24, 76, 312, 812, 1531, 6580) bytes => (6x6, 10x10, 15x15, 20x20, 40x40) matrix as shown in the Table 2.

Table 2. KRB22 encryption performance.

<i>File Size</i>	<i>ChaCha</i>	<i>BBRP22</i>
Two-Four	“1.69”	5.115
Seven-Six	“1.29”	5.742
Three-One- Two	“2.73”	8.994
Eight-Two- Two	“2.64”	11.119
One-Five- Three-One	“3.4”	10.939
Six-Five-Eight Zero	“2.27”	12.291

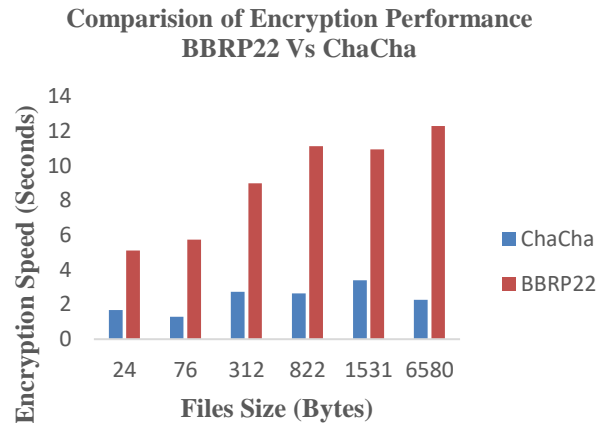


Figure 1. Encryption performance.

From Figure 1, the encryption execution assessment of the proposed calculation BBRP22 encryption execution contrasted and "ChaCha". "ChaCha" idea is the all askew qualities move to the first segment. The BBRP22 procedure has thought about the encryption speed in a moment or two. The encryption performance of the speed 5.115 (s), 5.742 (s), 8.994 (s), 11.119 (s), 10.939 (s) and 12.291 (s) for the" BBRP22. The BBRP22 gives more security of the information; when contrasted with existing procedures.

5. Conclusion

Now the world is fully activated with data alone. This data is produced by using online media; this data is unhindered data; this data does not have incredible security. To beat this issue, we will apply the ChaCha strategy. This method will effectively hack the information from the computer programmers. BBRP22 technique has 6 phases. To apply this technique, we must follow the given below steps: 1. To track down the great key S; 2. To get the Z1 & Z2 values from indivisible numbers; 3. To view as the an and b values in lattice; to find the values with the assistance k; and trade x and y values. 4. To get the T-test values and match those numbers. 5. To find the n esteem by utilizing k. 6. Apply the n values in legitimate condition; and trade x and y values. Hence, BBRP22 technique provides great security while comparing to ChaCha.

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